Mixing Time of Random Walk on Poisson Geometry Small World

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1 Small world effect

Small world effect, the fact that the diameters of most networks are considerably smaller than their sizes, is one of the most important features of real-world complex networks.

See graph G = (V, E) as a network, and suppose that |V| is large enough. We say G exhibits the small world effect, if the diameter of G is at most polynomially large in $\lg |V|$. Namely, for some polynomial function f, one has

$\operatorname{diam}(G) \le f(\lg |V|).$

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History

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1. Karinthy 1929, … a remarkable short story in his book: *Chains*.

2. Milgram 1960s, "six degrees of separation": Milgram carried out his famous "small-world" experiments, in which letters passed from person to person were able to reach a designated target individual within six steps.

3. Watts and Strogatz 1998, *Collective dynamics of 'small-world' networks*, Nature **393**, pp 440-442: WS small world.

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2 Why small world?—Modeling!

Three important models

1. BC small world 1988: B. Bollobás and F. Chung *The dimeter of a cycle plus a random matching*, SIAM J. Discrete Math. **1**, pp 328-333

2. WS small world 1998: D. J. Watts and S. H. Strogatz *Collective dynamics of 'small-world' networks*, Nature **393**, pp 440-442

3. NW small world 1999: M. E. J. Newman and D. J. Watts *Renormalization group analysis of the small-world network model*, Phys. Lett. A **263**, pp 341-346

Other models

LCD Model 2004: B. Bollobás and O. Riordan *The diameter of a scale-free random graph*, Combinatorica **24**, pp 5-34

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Remarks

1. 1-3 small world models revealed a common fact: that is, adding "long edges" to a regularly constructed (lattice-like) graph will make the resulted graph a small world—"adding long edges " mechanism.

The 1-dimensional lattice ring with n vertices is chosen to be the regularly constructed graph in all the three models.

2. Only the BC small world provided rigorous mathematical results.

3. A **NEW** mechanism other than the one working in BC, NW and WS small world makes the 'LCD' model a small world—we will discuss this aspect in another paper.

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3 Our model: Poisson geometry small world.

We consider the supercritical Poisson continuous percolation on d-dimensional torus T_n^d with volume n^d .

By adding "long edges" randomly to the largest percolation cluster, we obtain a random graph \mathcal{G}_n .

It can be proved that the diameter of \mathscr{G}_n grows at most polynomially fast in $\lg n$ and we call \mathscr{G}_n the *Poisson geometry small world*.

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Poisson continuous percolation on \mathbb{R}^d

Let \mathscr{P} denote the homogeneous Poisson point process of rate 1 on \mathbb{R}^d . Given r > 0. We define a random graph G = (V, E), where $V = \mathscr{P}$ and E consists of the edges which connect all Poisson point pairs lying in distance 2r.

Theorem A^{R.Meester and R.Roy (1995)} Suppose $d \ge 2$. Let $\psi(r)$ denote the probability that there exists a unique infinite connected component, write as C_{∞} , in G, let $\theta(r)$ denote the probability that the distance between the origin 0 and some vertex of C_{∞} is less than 2r. Then there exists $0 < r_c < \infty$ such that

$$\psi(r) \begin{cases} = 1, & \text{if } r > r_c \\ = 0, & \text{if } r < r_c \end{cases}, \quad \theta(r) \begin{cases} > 0, & \text{if } r > r_c \\ = 0, & \text{if } r < r_c \end{cases}$$

Where $\theta(r)$ is called the percolation probability.

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Continuous percolation on $B_n := [0, n]^d$

Let \mathscr{P}_n denote the homogeneous Poisson point process of rate 1 on B_n .

Given $r > r_c$. We define a random graph $G_n = (V_n, E_n)$, where $V_n = \mathscr{P}_n = B_n \cap \mathscr{P}$ and E_n consists of the edges which connect all \mathscr{P}_n point pairs lying in distance 2r.

A connected component C of G_n is called crossing for B_n , if the distance between C and any of the face of B_n is less than r.

For continuous percolation on $B_n := [0, n]^d$, we have the following two propositions.(Penrose and Pisztora (1996))

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Proposition 1 Suppose $r > r_c$. Suppose $\{\phi_n : n \ge 1\}$ is increasing with $\phi_n / \lg n \to \infty$ as $n \to \infty$, and with $\phi_n < n$ for all n.

Let $E_1(n)$ be the event that

(i) there is a unique component in G_n that is crossing for B_n , and (ii) no other component in G_n has diameter greater than ϕ_n . Then there exists a constant $c_1 > 0$ such that for all large enough n.

$$\mathbb{P}(E_1(n)) \ge 1 - \exp(-c_1\phi_n).$$

Remark. If we take $\phi_n = \lg^{(1+\epsilon)} n$, then it follows from Proposition 1 that, with high probability, all components except for the largest one have diameters less than $\lg^{(1+\epsilon)} n$.

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Proposition 2 Suppose $r > r_c$, and $0 < \epsilon < 1/2$. Let $E_2(n)$ be the event that

(i) there is a unique component C_b in G_n containing more than $\epsilon\theta(r)n^d$ points of \mathscr{P}_n ,

(ii)

$$(1-\epsilon)\theta(r) \le n^{-d}|C_b| \le (1+\epsilon)\theta(r),$$

(iii) C_b is crossing for B_n , and (iv) C_b is part of the infinite component C_{∞} in G. Then, there exist $c_2 > 0$ and n_2 such that

$$\mathbb{P}(E_2(n)) \ge 1 - \exp(-c_2 n^{d-1}), \quad n \ge n_2.$$

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Continuous percolation on the $d\mbox{-dimensional}$ torus T_n^d

Let T_n^d denote the *d*-dimensional torus obtained from B_n by cohering its opposite faces, and let \mathscr{P}_n^T denote the Poisson process of rate 1 on T_n^d .

Given $r > r_c$. We define a random graph $G_n^T = (V_n^T, E_n^T)$, where $V_n^T = \mathscr{P}_n^T$ and E_n^T consists of the edges which connect all \mathscr{P}_n^T point pairs lying in distance 2r.

Let C_{\max} denote the maximum connected component in G_n^T , let \bar{G}_n denote the subgraph of G_n^T , which corresponds the maximum connected component C_{\max} .

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Define the Poisson geometry small world \mathscr{G}_n

For any given constants α, β, σ and ζ satisfying $0 < \alpha < \beta < 1/2$, $\sigma > 0$ and $\zeta \ge 0$, we define a random graph $\mathscr{G}_n = \mathscr{G}_n(\alpha, \beta, \sigma, \zeta)$ from \overline{G}_n as the following: for any $u, v \in C_{\max}$, if

$$\alpha n \le d_{\infty}^T(u, v) \le \beta n,$$

then we connect u and v independently by a "long edge" with probability

$$p_n = \sigma n^{-d} \lg^{-\zeta} n;$$

otherwise, we do nothing.

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Where $d_{\infty}^T(\cdot, \cdot)$ denote the l_{∞} metric on T_n^d inherited from the usual l_{∞} metric $d_{\infty}(\cdot, \cdot)$ on \mathbb{R}^d defined by

$$d_{\infty}(x,y) := \max_{1 \le s \le d} |x_s - y_s|$$

for any $x, y \in \mathbb{R}^d$.

Let \mathcal{E}_n denote the new edge set with long edges, and let

 $\mathscr{G}_n = (C_{\max}, \mathcal{E}_n).$

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1. \mathscr{G}_n is a small world!

Theorem 1 Suppose $r > r_c$. Then for any $0 < \alpha < \beta < 1/2$ with

 $(2\beta)^d - (2\alpha)^d > 1/2,$

 $\sigma > 0$ and $\zeta \ge 0$, there exists constant $C_1 > 0$ such that for any $\nu > 0$,

 $\lim_{n \to \infty} \mathbb{P}(\operatorname{diam}(\mathscr{G}_n) \le C_1 \lg^{3+\zeta+\nu} n) = 1.$

Remark It seems that our setting on "long edge" is REASONABLE! Obviously, if only shorter edges, for example with length $n^{1-\epsilon}$, are added, then the diameter of the resulted graph grows at least fast in n^{ϵ} , and the resulted graph does not exhibit the small world effect. On the other hand, Theorem 1 indicates that, to make the resulted graph a small world, adding such shorter edges is not necessary.
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2. Random Walk on graph

In a graph G = (V, E), for any $u, v \in V$, let $d_G(u)$ be the the degree of u in G, and write $u \sim v$ if u and v are neighbors in G.

For any $u, v \in V$, we define a transition kernel by P(u, u) = 1/2, $P(u, v) = 1/2d_G(u)$ if $u \sim v$ and P(u, v) = 0 otherwise. A discrete time Markov chain $\{X_t : t \geq 0\}$ on V with transition kernel (P(u, v)) is called the *lazy random walk* on G.

Note that $\pi(u) := d_G(u)/D$ where $D = \sum_{v \in V} d_G(v) = 2|E|$, defines a reversible stationary distribution of $\{X_t\}$ since

$$\pi(u)P(u,v) = 1/2D = \pi(v)P(v,u).$$

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By the basic theory of Markov chains, for any initial state $u \in V$, the distribution of X_t , i.e. $P^t(u, \cdot) := \mathbb{P}(X_t \in \cdot \mid X_0 = u)$, converges weakly to π as $t \to \infty$. To measure convergence to equilibrium, we will use the *total variation distance*

$$||P^{t}(u, \cdot) - \pi||_{TV} := \sum_{v \in V} |p^{t}(u, v) - \pi(v)|.$$

The mixing time of $\{X_t : t \ge 0\}$ is defined by

$$T_{\min} := \min\left\{t : \max_{u \in V} ||P^t(u, \cdot) - \pi||_{TV} < 1/e\right\}.$$

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Random Walk on \mathscr{G}_n is rapid mixing

Denote by $T_{\rm mix}$ the mixing time of lazy random walk on $\mathscr{G}_n,$ then we have

Throrem 2 Suppose $r > r_c$. Then, (i) for any $0 < \alpha < \beta < 1/2$, $\zeta \ge d-1$ and for $\sigma > 0$ small enough, there exists constant $C_2 > 0$ such that

 $\lim_{n \to \infty} \mathbb{P}(T_{\min} \ge C_2 \lg n) = 1;$

(ii) for $0 < \alpha < \beta < 1/2$ with $(2\beta)^d - (2\alpha)^d > 1/2$, $\zeta \ge 0$ and $\sigma > 0$, there exists constant $C_3 > 0$ such that for any $\nu > 0$,

$$\lim_{n \to \infty} \mathbb{P}(T_{\text{mix}} \le C_3 \lg^{4\zeta + 5 + \nu} n) = 1.$$

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1. Small world effect

There is a known result in algebraic graph theory, see N. Alon and V. D. Milman (1985) and A. Berman and X. D. Zhang (2000), that

Lemma 1 For any connected graph G = (V, E), let $\Delta(G)$ denote its maximum degree, $\iota(G)$ denote its edge isoperimetric constant and let diam(G) denote its diameter. Then

diam(G)
$$\leq \frac{4\Delta(G)}{\iota(G)} \lg |V|.$$

The *edge isoperimetric constant* $\iota(G)$ is defined by

$$\iota(G) := \min_{S:|S| \le |V|/2} \frac{e(S, S^c)}{|S|}.$$

Where $e(S, S^c)$ be the number of edges between S and S^c .

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Theorem 1 follows from Lemma 1, Proposition 2 and the following Lemmas 2 and 3.

Lemma 2 Suppose $r > r_c$. Then

$$\lim_{n \to \infty} \mathbb{P}(\Delta(\mathscr{G}_n) \le \lg n) = 1.$$

Lemma 3 Suppose $r > r_c$, $0 < \alpha < \beta < 1/2$ with $(2\beta)^d - (2\alpha)^d > 1/2$, $\zeta \ge 0$ and $\sigma > 0$. Then there exists $C_4 > 0$ such that for any $\nu > 0$

$$\lim_{n \to \infty} \mathbb{P}\left(\iota(\mathscr{G}_n) \ge C_4 \lg^{-[1+\zeta+\nu]} n\right) = 1.$$

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2. Lower bound of the mixing time

To bound the mixing time from below, it suffices to prove that, with high probability (tends to 1 as $n \to \infty$), there exists a box B in T_n^d satisfying

1. The side length of B is $K \lg n$;

2. If we denote B' as the box with side length $K \lg n/2$ and centered at the same center of B. Then $B' \cap C_{\max} \neq \phi$. Note that C_{\max} is the vertex set of \mathscr{G}_n ;

3. There is no long edge between B and B^c .

Then the random walk started in B' can not escape from B in time $K \lg n/2r.$

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3. Upper bound of the mixing time

Conductance: For the lazy random walk $\{X_t : t \ge 0\}$ on \mathscr{G}_n , let $Q(u, v) := \pi(u)P(u, v)$ and $Q(S, S^c) := \sum_{u \in S} \sum_{v \in S^c} Q(u, v)$. Define

$$h := \min_{S:\pi(S) \le 1/2} \frac{Q(S, S^c)}{\pi(S)}$$

to be the *conductance* of $\{X_t : t \ge 0\}$. Letting $e(S, S^c)$ be the number of edges between S and S^c , we have

$$h = \frac{1}{2} \min_{S:\pi(S) \le 1/2} \frac{e(S, S^c)}{\text{Vol}(S)},$$

where $\operatorname{Vol}(S) = \sum_{u \in S} d_{\mathscr{G}_n}(u)$.

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For our lazy random walk $\{X_t : t \ge 0\}$ on \mathscr{G}_n , matrix theory tell us that the transition kernel (P(u,v)) has nonnegative real eigenvalues

$$1 = \lambda_0 \ge \lambda_1 \ge \lambda_2 \ge \ldots \ge \lambda_{|C_{\max}| - 1} \ge 0.$$

Note that $1 - \lambda_1$ is called the spectral gap of (P(u, v)). Let $\pi_{\min} = \min_{u \in C_{\max}} \pi(u)$.

It was given in A. Sinclair and M. Jerrum (1989) that

(*)
$$\sup_{u \in C_{\max}} ||P^t(u, \cdot) - \pi||_{TV} \le \frac{\lambda_1^t}{\pi_{\min}}.$$

On the other hand the spectral gap $1 - \lambda_1$ can be bounded from above and below by the conductance h in the following way (see Theorem 6.2.1] in R. Durrett (2007)),

$$(**) \qquad \frac{h^2}{2} \le 1 - \lambda_1 \le 2h.$$

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By equations (*) and (**) and the definition of mixing time, the upper bound for $T_{\rm mix}$ follows from the following lemma.

Lemma 4 Suppose $r > r_c$, $0 < \alpha < \beta < 1/2$ with

$$(2\beta)^d - (2\alpha)^d > 1/2,$$

 $\zeta \geq 0$ and $\sigma > 0$. Then there exists $C_5 > 0$ such that for any $\nu > 0$

$$\lim_{n \to \infty} \mathbb{P}\left(h \ge C_5 \lg^{-[2(\zeta+1)+\nu]} n\right) = 1.$$

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Geometry of \bar{G}_n

With high probability, we have

1).
$$(1 - \epsilon)\theta(r) \le n^{-d}|C_{\max}| \le (1 + \epsilon)\theta(r),$$

2). $(1 - \epsilon)\Gamma \le n^{-d}|\Lambda_n(u)| \le (1 + \epsilon)\Gamma$ for all $u \in C_{\max},$
3). $\Delta(\bar{G}_n) \le \lg n/2.$

Where $\Gamma = [(2eta)^d - (2lpha)^d] heta(r)$ and

$$\Lambda_n(u) = \left\{ v \in C_{\max} : \alpha n \le d_{\infty}^T(u, v) \le \beta n \right\}.$$

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The proof of Lemma 4 follows from the following two lemmas.

Lemma 5 Let $a = a(n) = \lg^{-(1+\zeta)} n$. Then

$$\lim_{n \to \infty} \mathbb{P}\left(\bigcap_{S:|S| \ge (1-a)|C_{\max}|} \left\{ \pi(S) > \frac{1}{2} \right\}\right) = 1.$$

Let

$$\hat{\iota}(\mathscr{G}_n) = \min_{S:|S| \le (1-a)|C_{\max}|} \frac{e(S, S^c)}{\operatorname{Vol}(S)}.$$

Then by Lemma 5,

$$\mathbb{P}\left(h \geq \hat{\iota}(\mathscr{G}_n)/2\right) \to 1, \text{ as } n \to \infty.$$

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So, to bound h from below, it suffices to bound $\hat{\iota}(\mathscr{G}_n)$ from bellow. In fact, we have

Lemma 6 Suppose $r > r_c$, $0 < \alpha < \beta < 1/2$ with

$$(2\beta)^d - (2\alpha)^d > 1/2,$$

 $\zeta \geq 0$ and $\sigma > 0$. Then there exists C > 0 such that for any $\nu > 0$

$$\lim_{n \to \infty} \mathbb{P}\left(\hat{\iota}(\mathscr{G}_n) \ge C \lg^{-[2(\zeta+1)+\nu]} n\right) = 1.$$

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The sketch of the proof of Lemma 6

Let

$$\mathcal{B}_{1} := \{ S \subset C_{\max} : |S| \le M/a \}, \text{ and} \\ \mathcal{B}_{2} := \{ S \subset C_{\max} : M/a < |S| \le a |C_{\max}| \}, \text{ and} \\ \mathcal{B}_{3} := \{ S \subset C_{\max} : a |C_{\max}| < |S| \le (1-a) |C_{\max}| \}, \end{cases}$$

where M > 0 is a large constant and $a = a(n) = \lg^{-(\zeta+1)} n$.

We will prove Lemma 6 for $S \in \mathcal{B}_i, i = 1, 2, 3$ respectively. Let

$$\mathcal{B} = \bigcup_{i=1}^{3} \mathcal{B}_i.$$

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For any
$$S \in \mathcal{B}$$
,

$$\frac{e(S, S^c)}{\operatorname{Vol}(S)} \ge \frac{c(S) + L(S, S^c)}{\Delta(\bar{G}_n)|S| + 2L(S)}$$

Where

c(S) is the number of connected components of S in \overline{G}_n . $L(S, S^c)$ is the number of "long edges" between S and S^c . L(S) is the number of "long edges" associated to S.

Clearly

$$L(S) = L(S, S^c) + L(S, S).$$

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Let $N(S, S^c)$ be the number of non-ordered pairs $u \in S, v \in S^c := C_{\max} \setminus S$ such that $u \in \Lambda_n(v)$, N(S, S) be the number of non-ordered pairs $u, v \in S$ such that $u \in \Lambda_n(v)$. Let

$$N(S) = N(S, S^c) + N(S, S).$$

Then $L(S, S^c)$ is independent to L(S, S), and

$$L(S, S^{c}) \sim b(N(S, S^{c}), p_{n}), \quad L(S, S) \sim b(N(S, S), p_{n}),$$

 $L(S) \sim b(N(S), p_n).$

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Large deviation inequality for binomial distribution.

Lemma 7: Suppose $Z \sim b(n, p)$. Then

 $\mathbb{P}(Z \geq zn) \leq \exp(-I(z)n), \; z > p; \; \mathbb{P}(Z \leq zn) \leq \exp(-I(z)n), \; z < p,$

where I(z) is the common rate function defined by

$$I(z) := z \lg \frac{zq}{(1-z)p} - \lg \frac{q}{1-z}, \ p \neq z \in (0,1).$$

Especially for small p, the above inequalities can be rewritten as

$$\mathbb{P}(Z \ge zpn) \le \exp(-\gamma(z)pn) \text{ for } z > 1, \text{ and}$$
$$\mathbb{P}(Z \le zpn) \le \exp\left(-\frac{1}{2}\gamma(z)pn\right) \text{ for } 0 < z < 1,$$

with $\gamma(z) = z \lg z - z + 1$.

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Now, let

$$\mathcal{B}_3 = \mathcal{B}_3^{>b} \cup \mathcal{B}_3^{$$

Where $\mathcal{B}_3^{>K}$ is the set of S in \mathcal{B}_3 such that $c(S) \ge |S|/\lg^b n$ and $\mathcal{B}_3^{<b} = \mathcal{B}_3 \setminus \mathcal{B}_3^{>b}$.

By the large deviation inequality, we have

$$\lim_{n \to \infty} \mathbb{P}_{\epsilon} \left(L(S) \le |S| \lg n, \ \forall \ S \in \mathcal{B}_3 \right) = 1.$$

Where \mathbb{P}_{ϵ} is the conditional probability function that the three items of the geometry of \bar{G}_n hold. This implies

$$\lim_{n \to \infty} \mathbb{P}_{\epsilon} \left\{ \frac{e(S, S^c)}{\operatorname{Vol}(S)} \ge \frac{2/3}{\lg^{b+1} n}, \ \forall \ S \in \mathcal{B}_3^{>b} \right\} = 1.$$

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Now, let's choose ϵ small enough such that

$$\frac{1-\epsilon}{(1+\epsilon)\tilde{\theta}(r)}\Gamma = \frac{(1-\epsilon)[(2\beta)^d - (2\alpha)^d]}{(1+\epsilon)} > 1/2.$$

For any $S\in \mathcal{B}_3$, if $a|C_{\max}|\leq |S|\leq |C_{\max}|/2$, then

$$\begin{split} N(S,S^c) &= \sum_{u \in S} \sum_{v \in S^c \cap \Lambda_n(u)} 1 \ge |S| \left(\min_{u \in S} |\Lambda_n(u)| - |S| \right) \\ &\ge |S| \left(\frac{1 - \epsilon}{(1 + \epsilon)\theta(r)} \Gamma |C_{\max}| - |S| \right) \\ &\ge a |C_{\max}| \left(\frac{1 - \epsilon}{(1 + \epsilon)\theta(r)} \Gamma |C_{\max}| - a |C_{\max}| \right) \\ &\ge a \left(\frac{1 - 2\epsilon}{(1 + \epsilon)\theta(r)} \Gamma \right) |C_{\max}|^2. \end{split}$$

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On the other hand, if $a|C_{\max}| \leq |S^c| \leq |C_{\max}|/2$, then

$$N(S, S^{c}) = N(S^{c}, S) = \sum_{u \in S^{c}} \sum_{v \in S \cap \Lambda_{n}(u)} 1 \ge |S^{c}| \left(\min_{u \in S^{c}} |\Lambda_{n}(u)| - |S^{c}| \right)$$
$$\ge a|C_{\max}| \left(\frac{1 - \epsilon}{(1 + \epsilon)\theta(r)} \Gamma|C_{\max}| - a|C_{\max}| \right)$$
$$\ge a \left(\frac{1 - 2\epsilon}{(1 + \epsilon)\theta(r)} \Gamma \right) |C_{\max}|^{2}.$$

So

$$N(S) \leq (1+\epsilon)\Gamma n^{d}|S| \leq (1+\epsilon)\Gamma n^{d}(1-a)|C_{\max}|$$

$$\leq \frac{(1+\epsilon)\Gamma(1-a)}{(1-\epsilon)\theta(r)}|C_{\max}|^{2} \leq \frac{(1+\epsilon)\Gamma}{(1-\epsilon)\theta(r)}|C_{\max}|^{2}$$

$$\leq \frac{1}{a} \cdot \frac{(1+\epsilon)^{2}}{(1-\epsilon)(1-2\epsilon)}N(S,S^{c}).$$

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Let
$$f(\epsilon) = (1-\epsilon)(1-2\epsilon)/(1+\epsilon)^2$$
, then

 $N(S,S^c) \geq f(\epsilon)aN(S).$

By the large deviation inequality, we have

$$\begin{aligned} \mathbb{P}_{\epsilon} \left(L(S, S^{c}) \geq f(\epsilon) a L(S) / 4 \right) \\ \geq \mathbb{P}_{\epsilon} (L(S, S^{c}) \geq \frac{1}{2} f(\epsilon) a N(S) p_{n} \text{ and } L(S) \leq 2N(S) p_{n}) \\ \geq 1 - \exp(-\frac{1}{4} \sigma \gamma(1/2) f(\epsilon) (1-\epsilon) \Gamma |S| \lg^{-(2\zeta+1)} n) \\ - \exp(-\frac{1}{2} \sigma \gamma(2) (1-\epsilon) \Gamma |S| \lg^{-\zeta} n) \\ \geq 1 - \exp\left(-C(\sigma, \epsilon) |S| \lg^{-(2\zeta+1)} n\right). \end{aligned}$$

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On the other hand, by the large deviation inequality, we have

$$\mathbb{P}_{\epsilon} \left(L(S) \ge \mu |S| \right) \ge \mathbb{P} \left(Z_2(S) \ge \frac{2\tau}{\sigma(1-\epsilon)\Gamma} N_2(S) p_n \right)$$
$$\ge 1 - \exp\left(-O\left(|S| \lg^{-\zeta} n \right) \right).$$

Where
$$\mu = \mu(n) = \tau \lg^{-\zeta} n$$
 with $2\tau/\sigma(1-\epsilon)\Gamma < 1$, $N_2(S) = \frac{1}{2}(1-\epsilon)\Gamma n^d |S|$ and

 $Z_2(S) \sim b(N_2(S), p_n).$

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Note that $L(S,S^c) \geq f(\epsilon) a L(S)/4$ and $L(S) \geq \mu |S|$ imply that

$$\frac{e(S,S^c)}{\operatorname{Vol}(S)} \quad \geq \frac{L(S,S^c)}{|S|\Delta(\bar{G}_n) + 2L(S)} \geq \frac{f(\epsilon)aL(S)/4}{|S|\lg n/2 + 2L(S)}$$

$$\geq \begin{cases} \frac{\tau \lg^{-\zeta} nf(\epsilon)a|S|/4}{|S|\lg n/2 + 2|S|} \geq \frac{C_6''}{\lg^{2(\zeta+1)} n}, & \text{if } L(S) \leq |S|; \\ \frac{f(\epsilon)aL(S)/4}{L(S)\lg n/2 + 2L(S)} \geq \frac{C_6'''}{\lg^{\zeta+2} n}, & \text{if } |S| \leq L(S). \end{cases}$$

Thus, for any $S \in \mathcal{B}_3$

$$\mathbb{P}_{\epsilon}\left(\frac{e(S,S^c)}{\operatorname{Vol}(S)} \ge \frac{C_6''}{\lg^{2(\zeta+1)}n}\right) \ge 1 - \exp\left(-C(\sigma,\epsilon)|S|\lg^{-(2\zeta+1)}n\right).$$

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Finally, for any $a|C_{\max}| \leq s \leq (1-a)|C_{\max}|$, let $\mathcal{B}_{3,s}^{< b}$ denote the set of $S \in \mathcal{B}_3^{< b}$ with |S| = s. For any $1 \leq j \leq s/\lg^b n =: j_s$, let $\mathcal{B}_{3,s,j}^{< b}$ denote the set of $S \in \mathcal{B}_{3,s}^{< b}$ such that c(S) = j. Then

$$\mathcal{B}_{3,s,j}^{\leq b}| \leq \binom{|C_{\max}|}{j} \binom{s-1}{j-1} \leq \binom{|C_{\max}|}{j} \binom{s}{j}$$

$$\leq \binom{|C_{\max}|}{j_s} \binom{s}{j_s} \leq \left(\frac{|C_{\max}|e}{j_s}\right)^{j_s} \cdot \left(\frac{se}{j_s}\right)^{j_s}$$
$$\leq \exp\left(j_s \left\{ \lg\left(\frac{|C_{\max}|}{s}\right) + 2\lg\left(\frac{s}{j_s}\right) + 2\right\} \right)$$

$$\leq \exp\left(D\lg\lg n\cdot s\lg^{-b}n\right).$$

Then

$$|\mathcal{B}_{3,s}^{\leq b}| \leq s \cdot \exp\left(D \lg \lg n \cdot s \lg^{-b} n\right).$$

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Provided $b = 2\zeta + 1 + \nu$ with $\nu > 0$, we have

$$\lim_{n \to \infty} \mathbb{P}_{\epsilon} \left(\frac{e(S, S^c)}{\operatorname{Vol}(S)} \ge \frac{C'_7}{\lg^{2(\zeta+1)} n}, \ \forall \ S \in \mathcal{B}_3^{< b} \right) = 1.$$

Thus, we finish the proof of the lemma.

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Thanks for your attention!

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